

# Symmetric Trivalent Categories

Josh Edge

Denison University

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# Planar diagrams

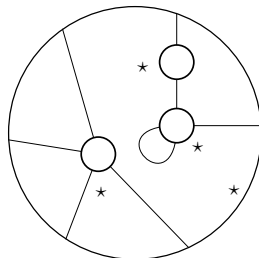
## Definition

A planar diagram consists of the following data:

- ① An outer disc and a finite number of inner discs.
- ② A finite number of marked boundary points on both outer and inner discs.
- ③ A collection of smooth, non-crossing curves called strings whose endpoints are the boundaries of the discs.
  - If there are  $2n$  total (inner and outer) boundary points, the planar diagram will have  $n$  strings connecting those boundary points.
  - The strings lie between the internal discs and external disc.
- ④ The outer disc and each inner disc will have a starred region.

# Example diagram

Diagrams are considered up to smooth isotopy. A planar diagram looks like

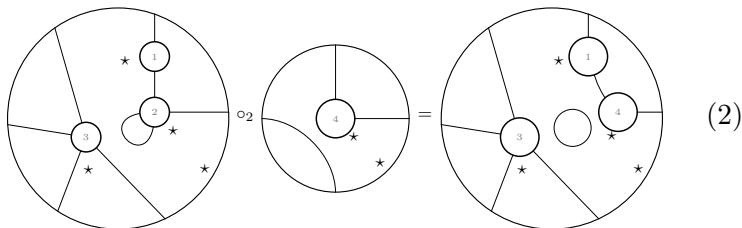


(1)

Note that these diagrams are sometimes called planar tangles, but this notation will be avoided due to the planar algebra we are working with.

# Composition of tangles

We can compose diagram  $A$  and  $B$  if the number of strands connecting to the outer boundary of  $B$  is equal to the number of strands connecting to one of the inner boundaries of  $A$ . An example of this composition is below:



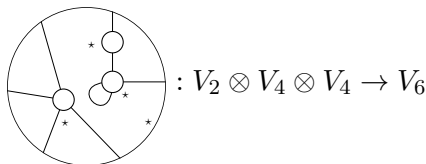
We call the set of isotopy classes of planar diagrams with the compositions described above the planar operad,  $\mathcal{P}$ .

# Planar Algebra Definition

## Definition

A planar algebra is a family of vector spaces  $\{V_k\}_{k=0,1,2,3,\dots}$  together with an action of the planar operad.

In (1), for example, a planar algebra would determine a multilinear map



that respects any composition like (2).

# Planar algebra definition

Some planar algebras will also have the following nice property:

## Definition

A planar algebra is said to have modulus  $d$  if inserting a closed string inside a planar diagram causes its multilinear map to be multiplied by  $d$ .

## Definition

A homomorphism of planar algebras is a collection of linear maps  $\{\phi_k : V_k \rightarrow W_k\}$  that intertwine the two actions of  $\mathcal{P}$ . We say that two planar algebras  $V$  and  $W$  are isomorphic if the  $\{\phi_k\}$  are isomorphisms.

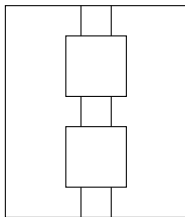
# Standard form

Often we draw planar diagrams in “standard form” as follows:

- ① Draw all boundaries (inner and outer) as rectangles parallel to the axes.
- ② Orient rectangle such that the sides have no strings and the left side of the rectangle is the starred region.
- ③ Typically, half of strings will go into the top of the diagram and half will go into the bottom.
- ④ All strings are drawn smoothly.

# Standard form example

- As an example consider the following diagram, which is commonly called the “multiplication” diagram



- Essentially, multiplying two diagrams in a planar algebra is the same as vertically stacking them.



# Skein Theories

- A natural class of examples of planar algebra comes from tangles as well as graphs.
- When one defines a ring or algebra, we often describe it using generators and relations.

## Definition

A planar algebra,  $\mathcal{V}$ , is generated by a set of elements,  $\mathcal{R}$ , if for every element  $v \in \mathcal{V}$ , there exists a planar diagram,  $T$ , such that  $T(r_1, \dots, r_k) = v$  for some  $r_i \in \mathcal{R}$ .

# Skein Theories

- When one studies all quotients of a ring or algebra, one might instead classify all two-sided ideals.

## Definition

Let  $\mathcal{V}$  be a planar algebra. A planar ideal,  $\mathcal{I}$ , is a subset of  $\mathcal{V}$  such that  $\mathcal{I}$  is the kernel of some quotient map of  $\mathcal{V}$ .

- Equivalently,  $\mathcal{I}$  is a planar ideal if it is closed under arbitrary composition with any element of the planar algebra.

## Definition

Let  $\mathcal{V}$  be a planar algebra. A *skein relation* is an element of the kernel of a quotient map from  $\mathcal{V}$ . A *skein theory* is a set of skein relations that generates a planar ideal.

- One of the most common examples of a planar algebra is the Temperley-Lieb-Jones planar algebra,  $\text{TLJ}(d)$ .
- It is the planar algebra generated by *no* objects.
- Thus, all elements of  $\text{TLJ}(d)$  contain only non-crossing strings. For example,

$$\text{crossing} = \text{two arcs} + 2 \cdot \text{vertical strand}$$

- The only relation is a circle relation:

$$\text{circle} = d$$

- In this way, one can think of any unshaded planar algebra with modulus  $d$  as containing  $\text{TLJ}(d)$  as a sub-planar algebra.
- The dimensions of the box spaces of  $\text{TLJ}(d)$  are indexed by the Catalan numbers  $(1, 1, 2, 5, \dots)$

- Another major class of planar algebras comes from those generated by



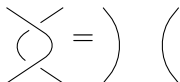
- With the following relations,  $\text{TLJ}(d)$  (with  $d = -(A^2 + A^{-2})$ ) can be thought of as a planar algebra generated by the crossing:

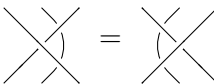
$$\bigcirc = -(A^2 + A^{-2})$$

$$\text{crossing with loop} = -A^{-3} \cup$$

$$\text{crossing} = A \left( + A^{-1} \cup \right)$$

- In addition,  $\text{TLJ}(d)$  satisfies the Reidemeister moves:

Reidemeister II (R2): 

Reidemeister III (R3): 

- Note that R1 is explicitly not satisfied!
- Obviously the box-space dimensions of  $\text{TLJ}(d)$  are the same  $(1, 1, 2, 5, 14 \dots)$ .

- Another common example of a planar algebra generated by a crossing is the Kauffman/Dubvronik polynomial planar algebra
- In addition to R2 and R3, it is subject to the following skein relations:

$$\begin{array}{c}
 \begin{array}{c} \diagup \diagdown \\ \diagdown \diagup \end{array} \pm \begin{array}{c} \diagdown \diagup \\ \diagup \diagdown \end{array} = z \left( \begin{array}{c} \left( \right) \end{array} \right) \left( \begin{array}{c} \pm \text{ (cup and cap) } \end{array} \right) \\
 \text{Circle} = d \qquad \text{Cup} = a \cdot \text{Cap}
 \end{array}$$

- Note that  $z$ ,  $d$ , and  $a$  are not all free parameters.
- The box-space dimensions are  $1, 1, 3, 15, \dots$

## Theorem

Let  $\mathcal{V}$  be a quotient of the planar algebra of tangles such that  $\dim V_4 = 2$  or  $3$ . Then there exists a relation of the form

$$\text{X} \pm \text{X} = z \left( \text{ } \right) \left( \pm \text{ } \right)$$

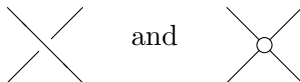
- Since  $\dim V_4 = 2$  or  $3$ , we know generically

$$a_1 \cdot \text{X} + a_2 \cdot \text{X} = a_3 \cdot \left( \text{ } \right) \left( + a_4 \cdot \text{ } \right)$$

- By rotating the equation however, we see that  $a_1 = \pm a_2$  and  $a_3 = \pm a_4$ . We also see that  $a_1 \neq 0$ .
- So we can rewrite the equation as

$$\text{X} \pm \text{X} = \frac{a_3}{a_1} \cdot \left( \text{ } \right) \left( \pm \text{ } \right)$$

- A consequence of this theorem is that  $\text{TLJ}(d)$  and the Kauffman/Dubrovnik polynomial planar algebras are the only planar algebras generated by the crossing with  $\dim V_4 \leq 3$ .
- It is this type of classification of “small” planar algebras that we will be concerned with today.
- One might also consider what small planar algebras exist up to isomorphism that are generated by



- This second crossing is called a virtual crossing.
- Can think of it as literally passing “through” the strands.



## Theorem (E. '18)

*The only small planar algebras generated by a crossing and a virtual crossing are isomorphic to one of the following:*

- ① *A virtual TLJ planar algebra*
- ② *The Kauffman polynomial planar algebra at  $a = \pm 1$  and  $d = -2$  equipped with the virtual crossing*

$$\text{Virtual Crossing} = - \left( \text{Cap} \right) \left( + \text{Cup} \right)$$

- ③  *$\text{Rep}(O(2), a)$ , where  $a \in \mathbb{C}^\times$  denotes that the planar algebra is equipped with a braiding given by the formula*

$$\text{Crossing} = \frac{a^{-1} - a}{2} \left( \text{Cap} \right) \left( - \frac{a^{-1} - a}{2} \text{Cup} + \frac{a^{-1} + a}{2} \text{Virtual Crossing} \right) \quad (3)$$

- Another class of planar algebras are those of graphs.
- In our case, we will consider graphs generated by a symmetric, trivalent vertex.



- Elements of this planar algebra can be thought of as trivalent graphs, graphs with only trivalent and univalent vertices
- Because strings cannot cross, these graphs must also be *planar*

# Previous work

In 2017 work of Morrison, Peters, and Snyder, all planar algebras generated by



were classified with certain smallness conditions and subject to the following relations:

$$= d = t - 1$$

$$= 0$$

$$= c_1 \cdot$$

$$= c_2 \cdot$$

# Previous work

## Theorem (From MPS 2017)

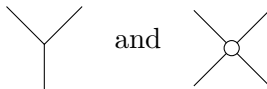
*The following table is an exhaustive list of non-degenerate trivalent planar algebras with initial box space dimension bounds 1, 0, 1, 1, 4, 11, 41:*

<i>Dimension bounds</i>	<i>Name</i>
1, 0, 1, 1, 2, ...	$SO(3)_{\zeta_5}$
1, 0, 1, 1, 3, ...	$SO(3)_q$ or $OSp(1 2)$
1, 0, 1, 1, 4, 8, ...	$ABA$
1, 0, 1, 1, 4, 9, ...	$(G_2)_{\zeta_{20}}$
1, 0, 1, 1, 4, 10, ...	$(G_2)_q$
1, 0, 1, 1, 4, 11, 37, ...	$H3$

*where  $ABA$  is a sub-planar algebra of the free product of  $TL(\sqrt{dt^{-1}}) * TL(t)$  and  $H3$  is the fusion category found in work by Grossman and Snyder 2012.*

# Symmetric Trivalent Categories

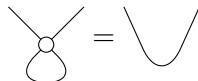
My work expands on this classification by looking at planar algebras generated by



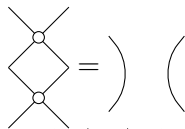
# Virtual Reidemeister Moves

The virtual crossing has corresponding virtual Reidemeister moves:

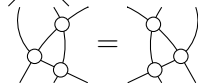
Virtual Reidemeister I (vR1):



Virtual Reidemeister II (vR2):



Virtual Reidemeister III (vR3):



in addition to some naturality conditions with the trivalent vertex.

*Equivalently, one can think of adding the virtual crossing as simply removing the planarity conditions from the graphs.*

# Computer Use

Because the virtual crossing allows us to consider any graph, we can more easily utilize Mathematica to answer our questions.

```
in[4]:= STtoGraph[Collect[R3L, listR32]]
```

$$\begin{aligned}
 \text{Out[4]} = & \left( x^3 + x z^2 \right) \begin{array}{|c|} \hline \text{Diagram 1} \\ \hline \end{array} + x^2 y \begin{array}{|c|} \hline \text{Diagram 2} \\ \hline \end{array} + x^2 z \begin{array}{|c|} \hline \text{Diagram 3} \\ \hline \end{array} + \left( -\frac{wxy}{-2+t} + xy^2 + yz^2 \right) \begin{array}{|c|} \hline \text{Diagram 4} \\ \hline \end{array} + \left( -\frac{wyz}{-2+t} + xyz + y^2 z \right) \begin{array}{|c|} \hline \text{Diagram 5} \\ \hline \end{array} + \\
 & \left( \frac{2wxy}{-2+t} + 2x^2 y + wy^2 - xy^2 + tx y^2 + y^3 + \frac{2wyz}{-2+t} + 2xyz + y^2 z \right) \begin{array}{|c|} \hline \text{Diagram 6} \\ \hline \end{array} + 2x^2 z \begin{array}{|c|} \hline \text{Diagram 7} \\ \hline \end{array} + xyz \begin{array}{|c|} \hline \text{Diagram 8} \\ \hline \end{array} + xz^2 \begin{array}{|c|} \hline \text{Diagram 9} \\ \hline \end{array} + \\
 & xyz \begin{array}{|c|} \hline \text{Diagram 10} \\ \hline \end{array} + yz^2 \begin{array}{|c|} \hline \text{Diagram 11} \\ \hline \end{array} + \left( -\frac{wyz}{-2+t} + xyz + y^2 z \right) \begin{array}{|c|} \hline \text{Diagram 12} \\ \hline \end{array} + xz^2 \begin{array}{|c|} \hline \text{Diagram 13} \\ \hline \end{array} - z^3 \begin{array}{|c|} \hline \text{Diagram 14} \\ \hline \end{array} + \left( -\frac{wxy}{-2+t} - xy^2 + yz^2 \right) \begin{array}{|c|} \hline \text{Diagram 15} \\ \hline \end{array} - \\
 & w^2 x \begin{array}{|c|} \hline \text{Diagram 16} \\ \hline \end{array} - w^2 x \begin{array}{|c|} \hline \text{Diagram 17} \\ \hline \end{array} - w^2 z \begin{array}{|c|} \hline \text{Diagram 18} \\ \hline \end{array} + \left( \frac{w^3}{(-2+t)^2} - \frac{w^2 y}{-2+t} \right) \begin{array}{|c|} \hline \text{Diagram 19} \\ \hline \end{array} - w^2 z \begin{array}{|c|} \hline \text{Diagram 20} \\ \hline \end{array} - w^2 z \begin{array}{|c|} \hline \text{Diagram 21} \\ \hline \end{array} +
 \end{aligned}$$

# Simple skein theory classification

Why should we care?

- Small objects frequently show up in our calculations, especially in classification questions.
- Having an ability to discern which small object you are working with quickly is very useful.
- In general, classification problems can help us find objects we didn't know existed.



# Simple skein theory classification

Some of the planar algebras in MPS17 will also be naturally included in this classification. For instance  $\text{OSp}(1|2)$  has relations:


$$\begin{array}{c} \diagup \quad \diagdown \\ \circ \\ \diagdown \quad \diagup \end{array} = -2 \cdot \left( \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right) \quad \left( \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right) = - \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} + 2 \cdot \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array}$$

$$\begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} - \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} = -\frac{1}{2} \left[ \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right] \quad \left( \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right) = - \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array}$$

# Simple skein theory classification

- Of course there is no point in doing this classification if nothing new exists!
- One example not included in MPS17 is Deligne's  $S_t$ , which has I=H relation:

$$\begin{array}{c} \diagup \quad \diagdown \\ | \\ \diagdown \quad \diagup \end{array} - \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \end{array} = \frac{1}{t-2} \left( \begin{array}{c} \diagup \quad \diagdown \\ | \\ \diagdown \quad \diagup \end{array} \right) - \left( \begin{array}{c} \diagdown \quad \diagup \\ | \\ \diagup \quad \diagdown \end{array} \right)$$

-  cannot be written as a linear combination of non-virtual diagrams.
- At values of  $t \in \mathbb{Z}_{>0}$ , Deligne's  $S_t$  is a degenerate version of  $\text{Rep}(S_t)$ .

# Classification

## Theorem

Consider  $\mathcal{T} = \{T_n\}_{n=0,1,2,\dots}$ , a planar algebra generated by a trivalent vertex and virtual crossing. Any non-fully flat quotient of  $\mathcal{T}$  with dimension bounds “1,0,1,1,4,...” is one of the following:

Dimensions	Name
1, 0, 1, 1, 3, ...	$Rep(SO(3))$ , $Rep(OSp(1 2))$ , and $Rep(S_3)$
1, 0, 1, 1, 4, 9, ...	$Rep(G_2)$
1, 0, 1, 1, 4, 10, ...	$Rep(S_4)$
1, 0, 1, 1, 4, 11, ...	$S_t$ for $t \notin \mathbb{Z}_{\geq 0}$
1, 0, 1, 1, 4, 11, ...	$Rep(S_t)$ for $t \in \mathbb{Z}_{\geq 0} - \{0, 1, 2, 3\}$

## Lemma

*Suppose we have a symmetric trivalent planar algebra with  $t \neq 0$  or 2. Then a planar algebra has a relation of the form*

$$\begin{array}{c} \diagup \quad \diagdown \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} + \begin{array}{c} \diagdown \quad \diagup \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagup \quad \diagdown \end{array} = z \cdot \left[ \begin{array}{c} \diagup \quad \diagdown \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right] \left( + \begin{array}{c} \diagdown \quad \diagup \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagup \quad \diagdown \end{array} \right)$$

*if and only if it has a relation of the form*

$$\begin{array}{c} \diagup \quad \diagdown \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} - \begin{array}{c} \diagdown \quad \diagup \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagup \quad \diagdown \end{array} = z' \cdot \left[ \begin{array}{c} \diagup \quad \diagdown \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array} \right] \left( - \begin{array}{c} \diagdown \quad \diagup \\ \quad \quad \diagdown \quad \diagup \\ \quad \quad \diagup \quad \diagdown \\ \diagup \quad \diagdown \end{array} \right)$$

*with  $z = \frac{1}{t}$  and  $z' = \frac{1}{t-2}$ .*

- Multiplying the top equation by an H we get

$$\begin{array}{c} \diagup \quad \diagdown \\ \square \\ \diagdown \quad \diagup \end{array} = \frac{t-1}{t} \cdot \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \end{array} + \frac{1}{t} \cdot \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagup \quad \diagdown \end{array} + \frac{1}{t} \cdot \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \end{array}$$

- The left side is rotationally invariant, so is the right side!
- Rotating and setting the result equal to each other gives us

$$\begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \end{array} - \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagup \quad \diagdown \end{array} = \frac{1}{t-2} \cdot \left[ \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \end{array} \quad \begin{array}{c} \diagup \quad \diagdown \\ \text{---} \\ \diagdown \quad \diagup \end{array} \right]$$

# Classification

- Unfortunately (fortunately?), most of the other calculations were done using Mathematica.
- You can download my code from Arxiv and play around with it, but you may have to download the Knot package.
- Essentially, the only symmetric trivalent planar algebras with these conditions are already enumerated in MPS17 or Deligne's  $S_t$ .
- The question of which of these categories is braided is also interesting.

# Sub-braidings

## Theorem

*The only non-fully-flat sub-braidings of  $S_t$  for generic  $t$  and  $\text{Rep}(S_t)$  for  $t \in \mathbb{Z}_{\geq 0} - \{0, 1, 2, 3\}$  is the sub-braiding inherited from  $SO(3)_q$ :*

$$\text{Virtual Crossing} = (q^2 - 1) \cdot \left( \text{Cap} + q^{-2} \text{Cup} - (q^2 + q^{-2}) \text{Trivalent} \right)$$

*for  $t = q^2 + 2 + q^{-2}$ . At  $t = 0$  and  $3$  all sub-braidings are fully-flat.*

- This is not a true braiding on the category—no naturality condition with the virtual crossing.
- There is a map from  $SO(3)_q \hookrightarrow \text{Rep}(S_t)$  when  $t = q^2 + 2 + q^{-2}$ .
- Interesting that there are no other sub-braidings.

## Future Work

- Repeating this work in characteristic  $p$ .
- Working with directed graphs.
- Both are great projects for undergraduates!



Thank you!

Thank you for listening!